

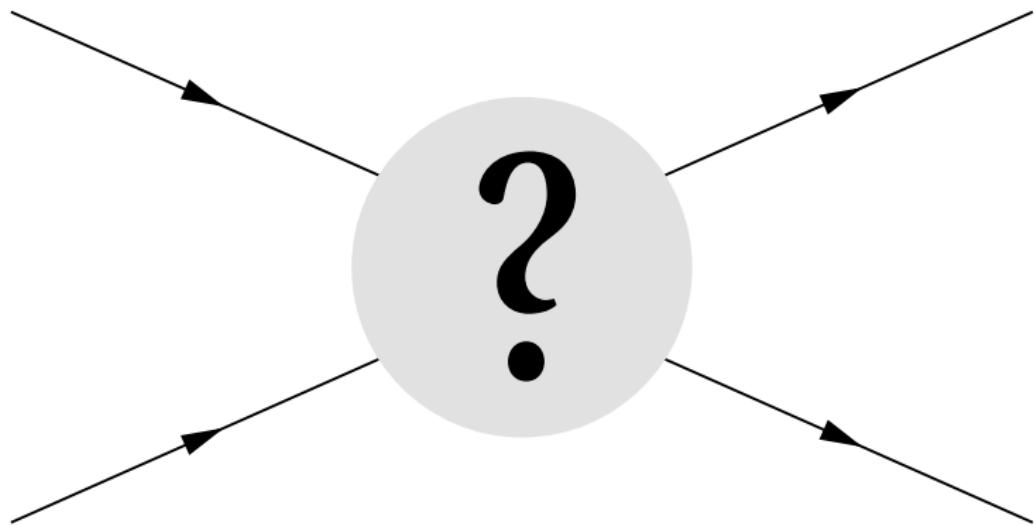
# Parametric Feynman integrals with hyperlogarithms

Erik Panzer

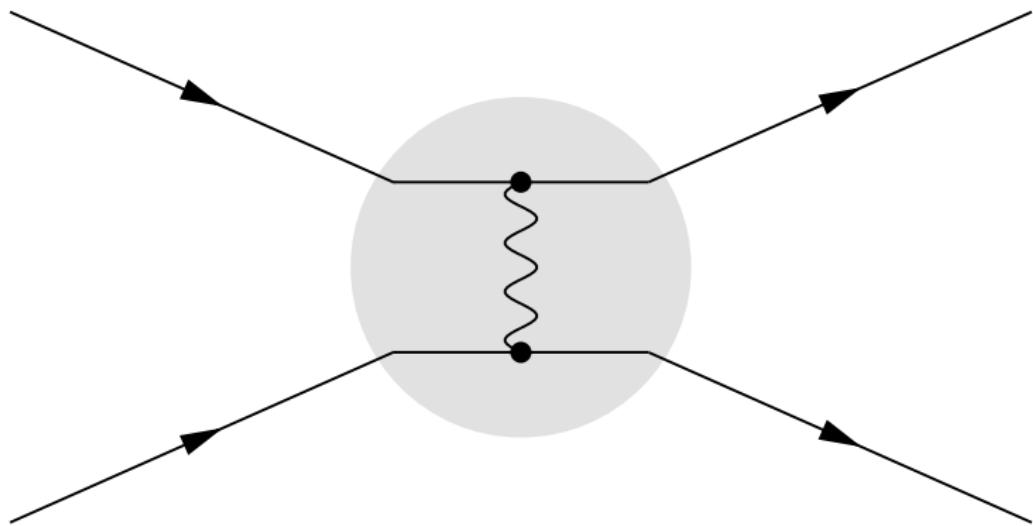
All Souls College

February 15th  
Trinity College Dublin

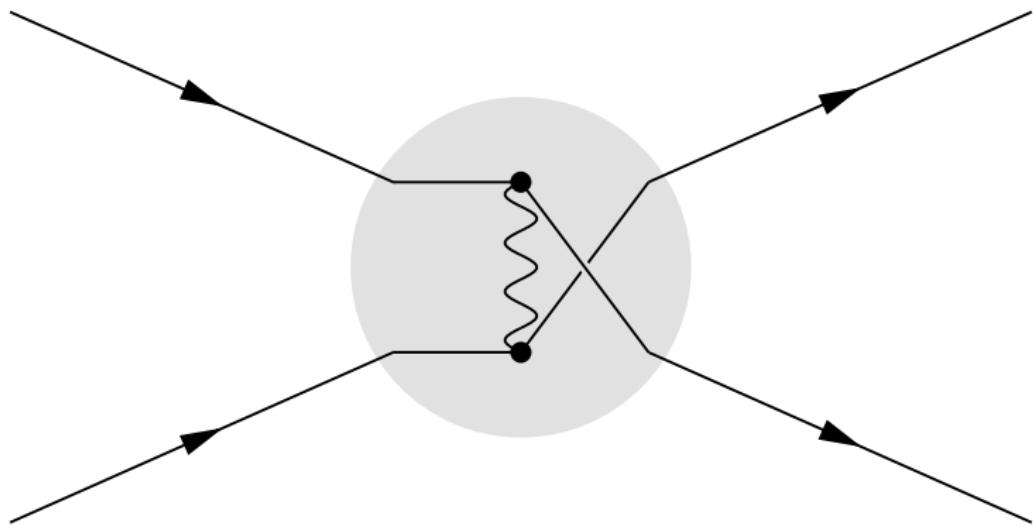
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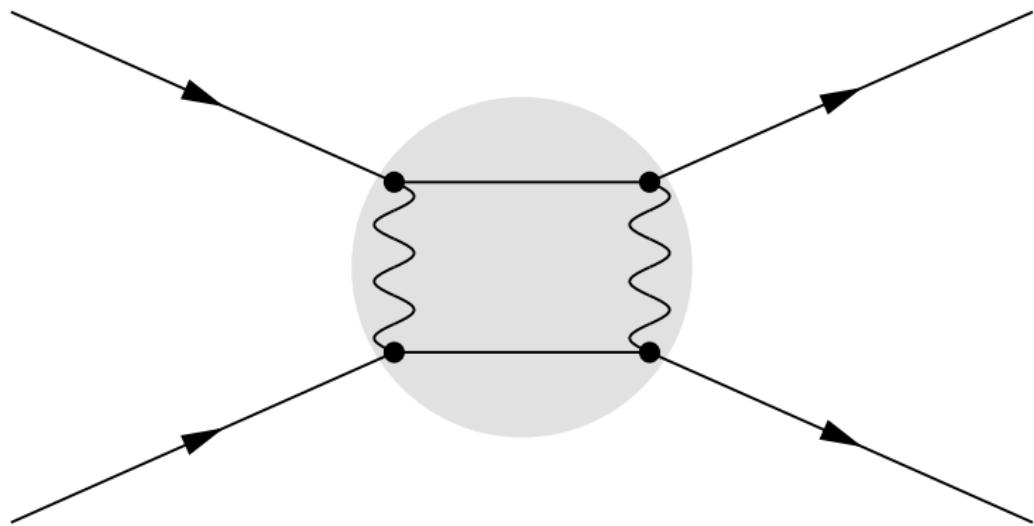
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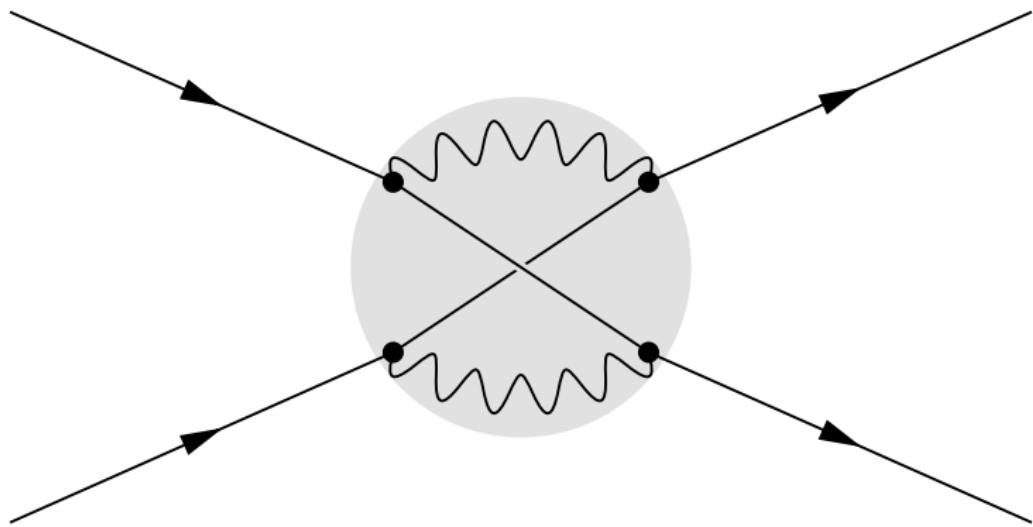
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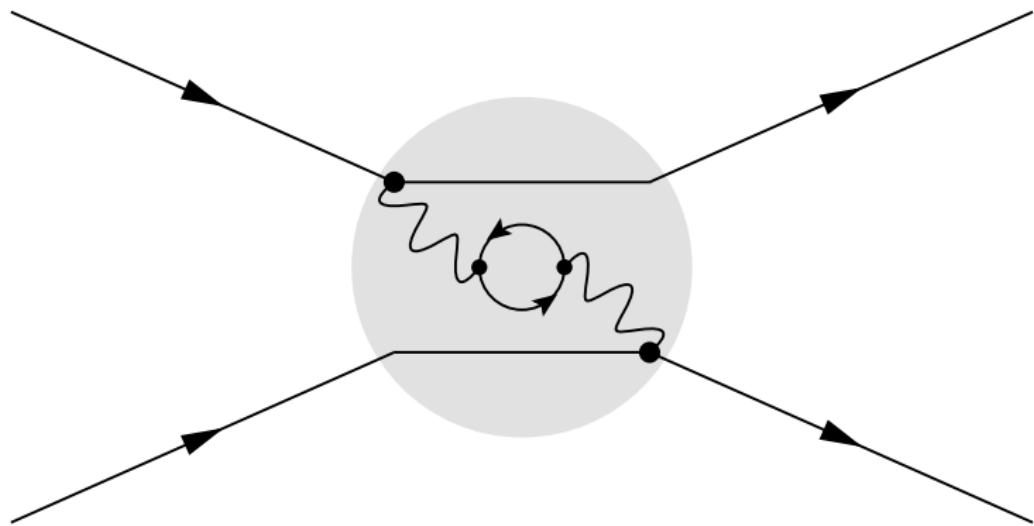
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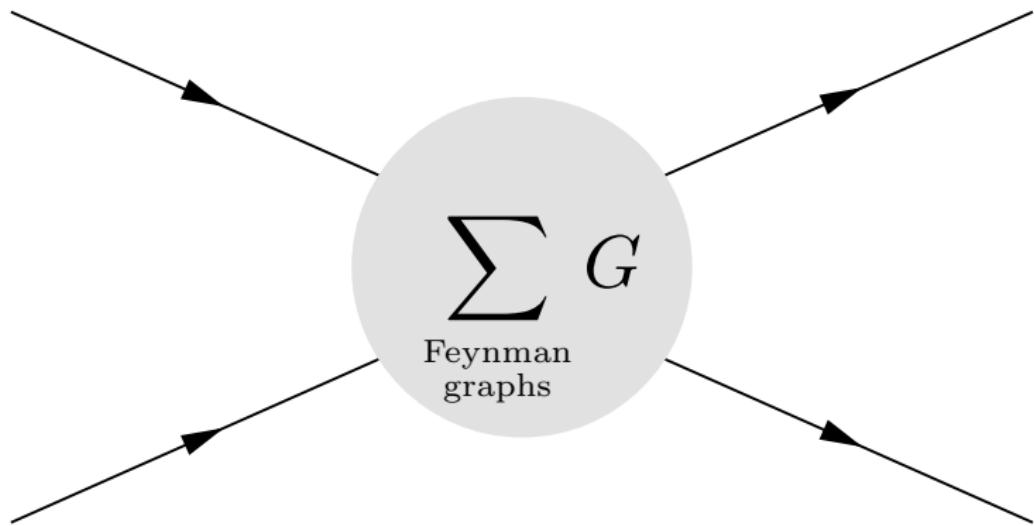
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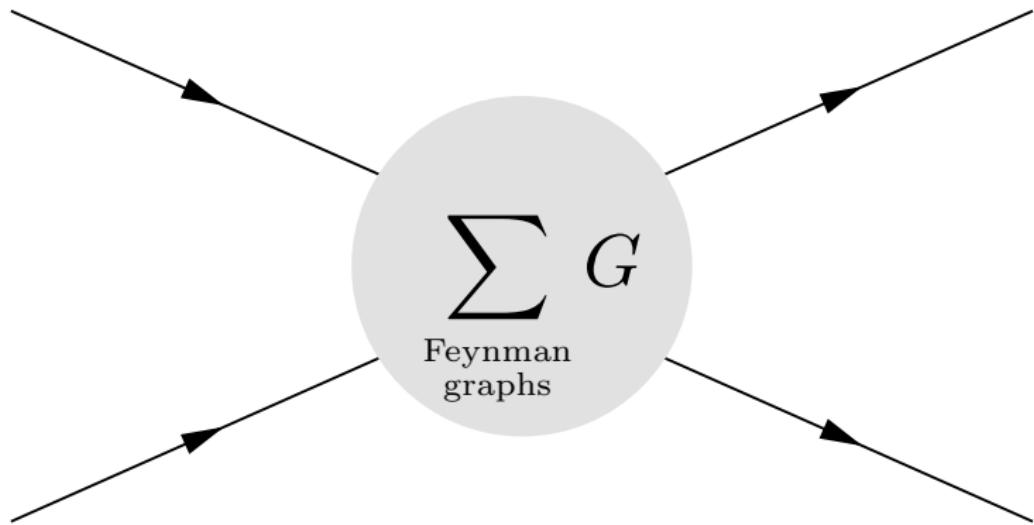
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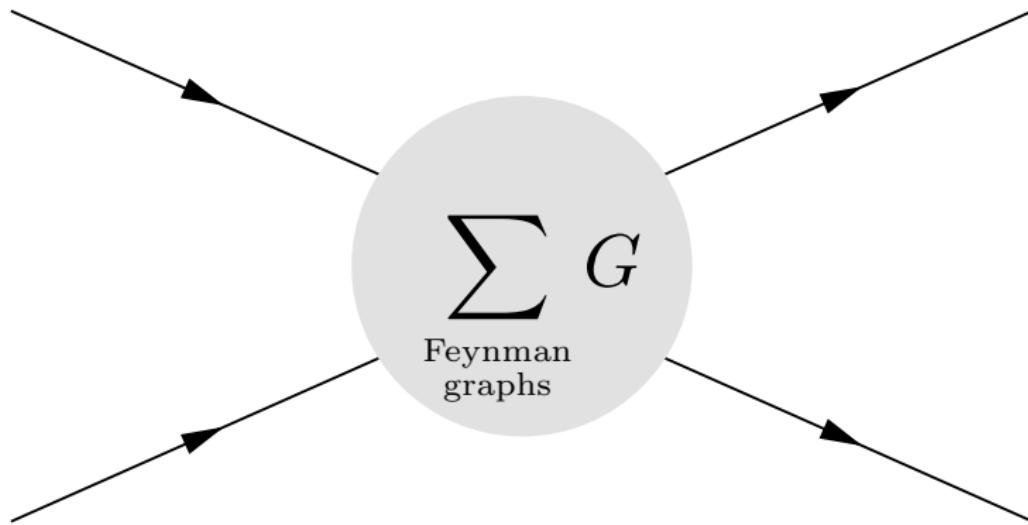


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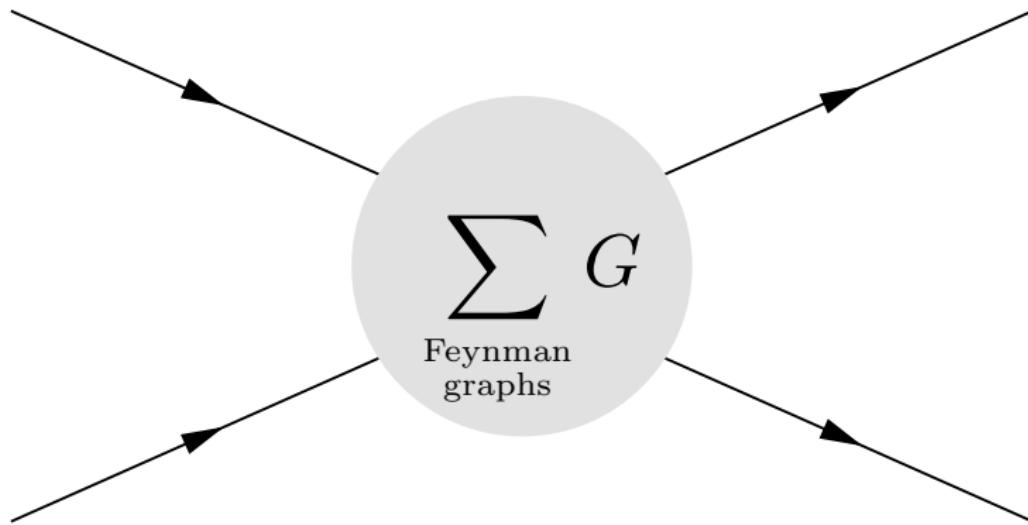
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Some FI are expressible with

- logarithms:

$$-\log(1 - z) = \sum_{0 < k} \frac{z^k}{k}$$

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## Schwinger parameters

With the *superficial degree of divergence*  $\text{sdd} = |E(G)| - D/2 \cdot \text{loops}(G)$ ,

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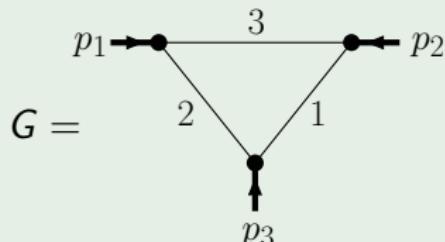
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Example ( $D = 4$ )



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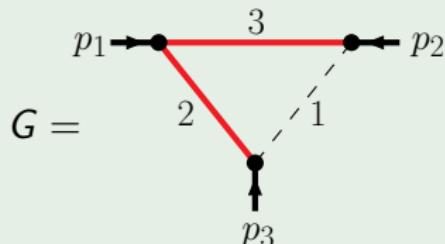
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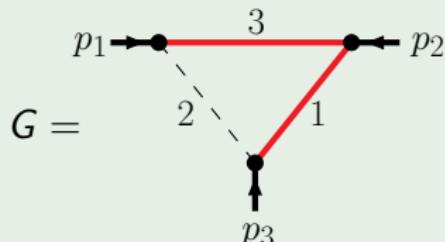
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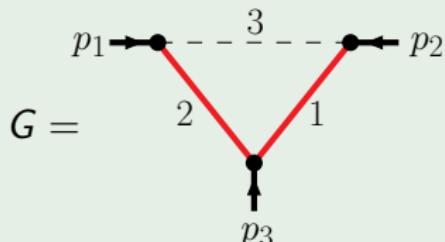
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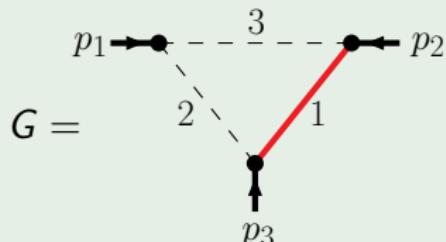
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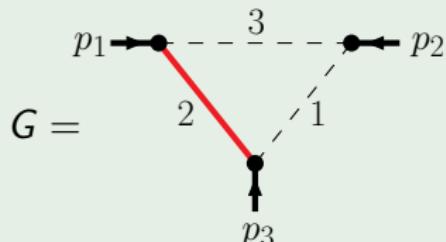
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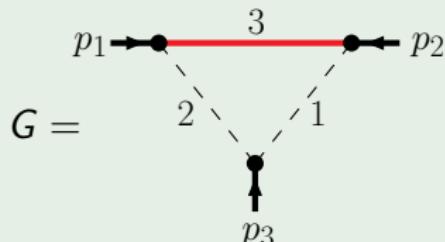
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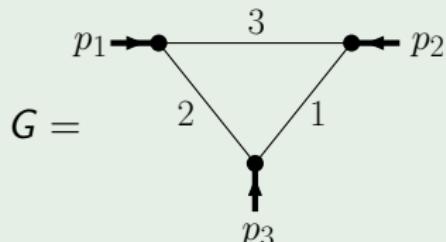
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$$\{t_3\} \times \{t_1 \leq t_2\} = \{t_1 \leq t_2 \leq t_3\} \cup \{t_1 \leq t_3 \leq t_2\} \cup \{t_3 \leq t_1 \leq t_2\}$$

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- multivalued, monodromies, path concatenation
- represent all MPL:

$$G(0^{n_d-1}, \sigma_d, \dots, 0^{n_1-1}, \sigma_1; z) = (-1)^d \text{Li}_{n_1, \dots, n_d} \left( \frac{\sigma_2}{\sigma_1}, \dots, \frac{\sigma_d}{\sigma_{d-1}}, \frac{z}{\sigma_d} \right)$$

# Symbolic integration: Rewriting hyperlogarithms

Problem: Given  $G(\vec{\sigma}(\alpha); z)$ , write it as hyperlogarithm with constant letters and final argument  $\alpha$ . Recursive solution via

$$dG(\vec{\sigma}; z) = \sum_{i=1}^n G(\dots, \phi_i, \dots; z) d \log \frac{\sigma_i - \sigma_{i-1}}{\sigma_i - \sigma_{i+1}} \quad \sigma_0 := z, \sigma_{n+1} := 0$$

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Computation of integration constants relies on shuffle algebra, rescalings

$$G(\lambda \vec{\sigma}; \lambda z) = G(\vec{\sigma}; z)$$

and Möbius transformations.

All this is implemented in the Maple program `HyperInt`.

## Linear reducibility

We need that all partial integrals

$$f_n := \int_0^\infty f_{n-1} \, d\alpha_n = \int_{(0,\infty)^n} f_0 \, d\alpha_1 \cdots d\alpha_n \quad \left( f_0 = \frac{1}{\psi^{D/2-\text{sdd}} \varphi^{\text{sdd}}} \right)$$

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- condition on the polynomials  $\psi$  and  $\varphi$  only
- sufficient criteria: polynomial reduction algorithms (Brown)

# Polynomial reduction

Denote alphabets (divisors) by sets  $S$  of irreducible polynomials.

## Definition

Let  $S$  denote a set of polynomials  $f = f^e \alpha_e + f_e$  linear in  $\alpha_e$ . Then with  $[f, g]_e := f^e g_e - f_e g^e$ ,  $S_e$  shall be the set of irreducible factors of

$$\{f^e, f_e : f \in S\} \quad \text{and} \quad \{[f, g]_e : f, g \in S\}.$$

## Example (massless triangle)

$$S = \{\psi, \varphi\} = \{\alpha_1 + \alpha_2 + \alpha_3, \alpha_2 \alpha_3 + z \bar{z} \alpha_1 \alpha_3 + (1-z)(1-\bar{z}) \alpha_1 \alpha_2\}$$
$$[\varphi, \psi]_3 = (\alpha_1 + \alpha_2)(\alpha_2 + \alpha_1 z \bar{z}) - (1-z)(1-\bar{z}) \alpha_1 \alpha_2$$

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$$S_3 = \{\alpha_1 + \alpha_2, z \alpha_1 + \alpha_2, \bar{z} \alpha_1 + \alpha_2, z \bar{z} \alpha_1 + \alpha_2, \alpha_1, \alpha_2, 1-z, 1-\bar{z}\}$$

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## Lemma

If the singularities of  $F$  are contained in  $S$ , then the singularities of  $\int_0^\infty F d\alpha_e$  are contained in  $S_e$ .

# Polynomial reduction

## Corollary (linear reducibility)

If all  $S^k := (S^{k-1})_k$  are linear in  $\alpha_{k+1}$ , then any MPL  $F$  with alphabet in  $S^0$  integrates to a MPL  $\int_0^\infty F \prod_{e=1}^n d\alpha_e$  with alphabet in  $S^n$ .

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This gives only very coarse upper bounds, for example  $z\bar{z}-1$  is spurious:  
It drops out in  $S_{2,3} \cap S_{3,2} = \{z, \bar{z}, 1-z, 1-\bar{z}, z-\bar{z}\}$  because

$$S_{2,3} = \{z, \bar{z}, 1-z, 1-\bar{z}, z-\bar{z}, z\bar{z}-z-\bar{z}\}.$$

Note that  $z\bar{z}-z-\bar{z}$  is spurious.

# Compatibility graphs

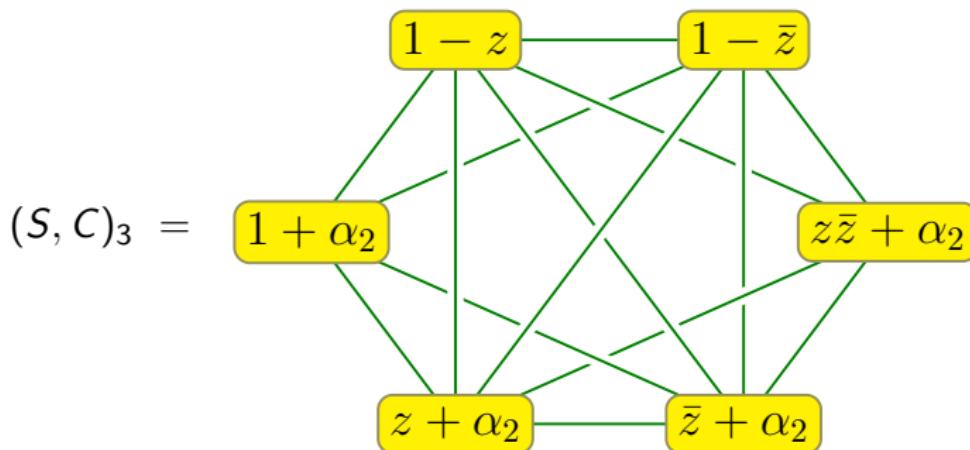
Keep track of **compatibilities**  $C \subset {S \choose 2}$  between polynomials:

- start with the complete graph  $\psi$  —  $\varphi$
- in  $S_e$ , only take resultants  $[f, g]_e$  for compatible  $\{f, g\} \in C$
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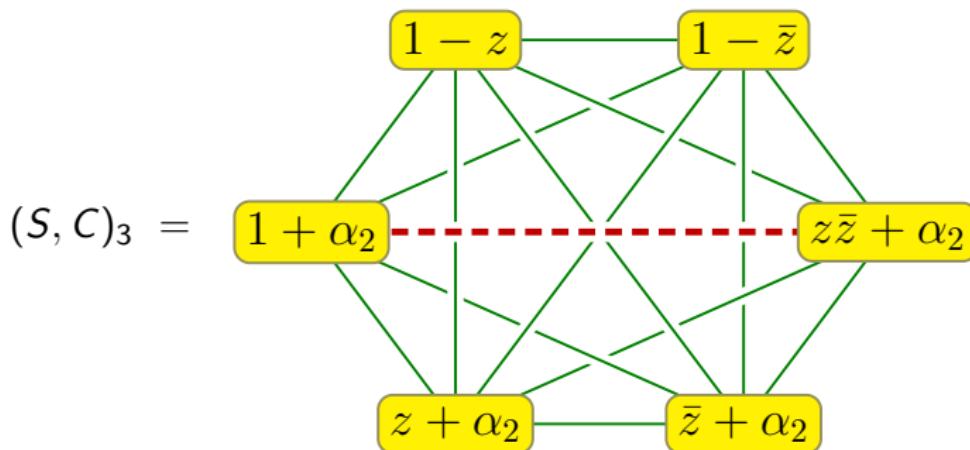
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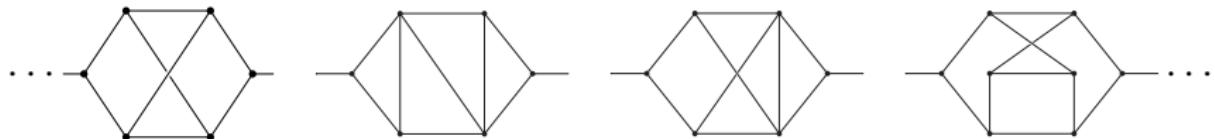
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$z\bar{z}\alpha_1 + \alpha_2$  and  $\alpha_1 + \alpha_2$  not compatible  $\Rightarrow$  no resultant  $1 - z\bar{z}$  in  $(S, C)_{3,2}$

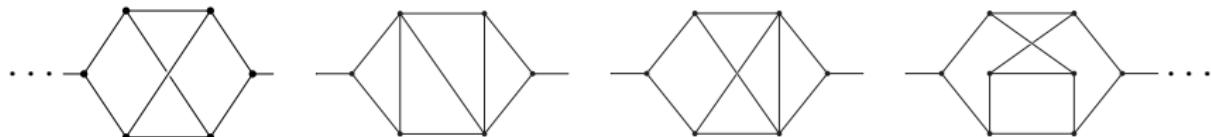
# Linearly reducible families (fixed loop order)

- ① all  $\leq 4$  loop massless propagators (Panzer)

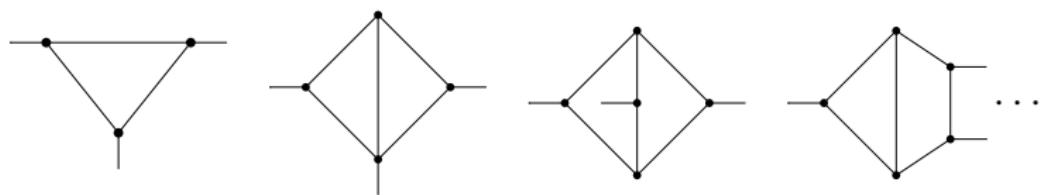


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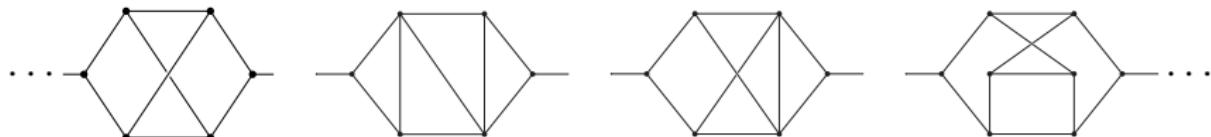


- ② all  $\leq 3$  loop massless off-shell 3-point (Chavez & Duhr, Panzer)

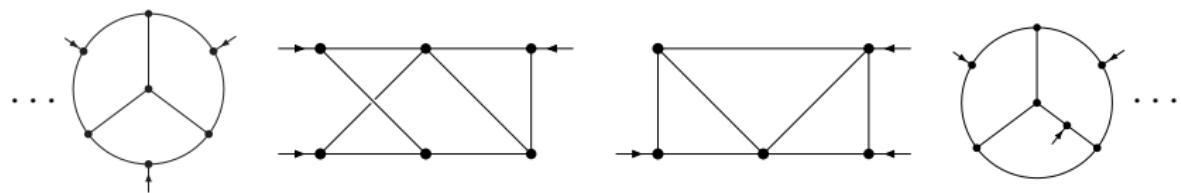


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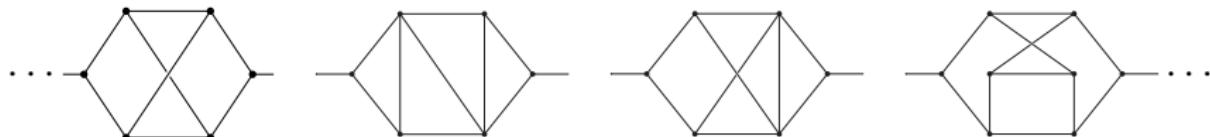


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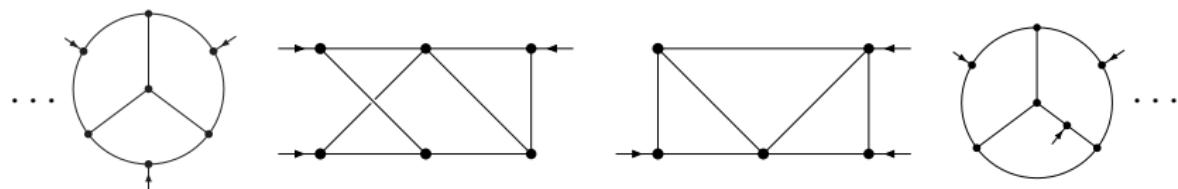


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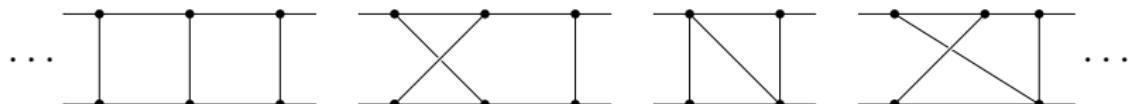
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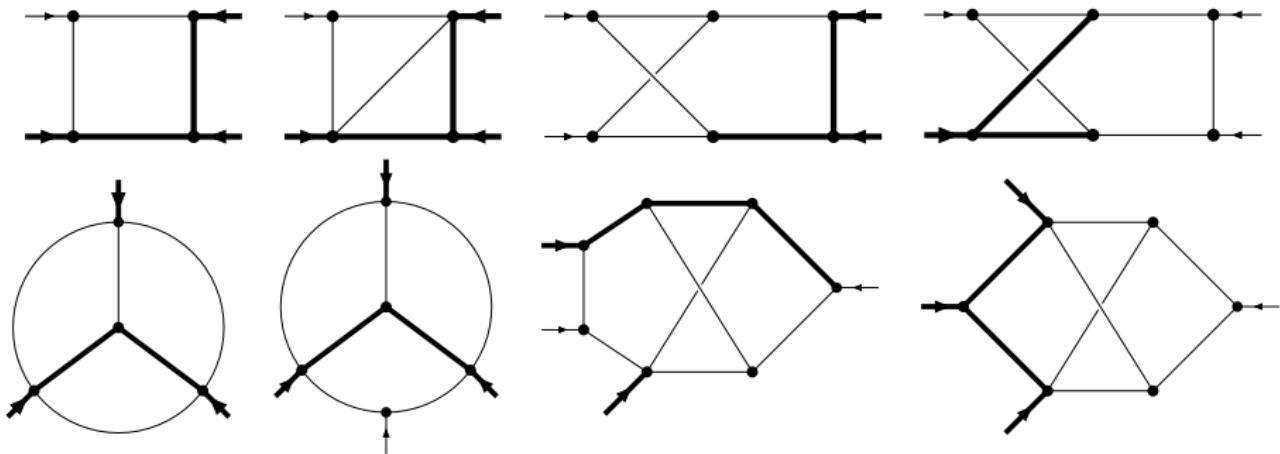
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- ③ all  $\leq 2$  loop massless on-shell 4-point (Lüders)



## Linearly reducible massive graphs (some examples)

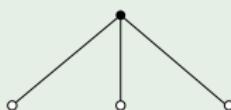


# Linearly reducible families (infinite)

- 3-constructible graphs (3-point functions) [Brown, Schnetz, Panzer]



## Example

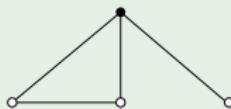


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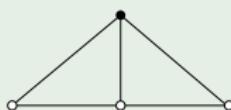


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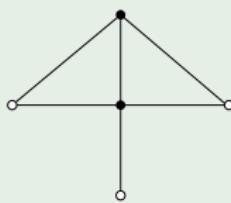


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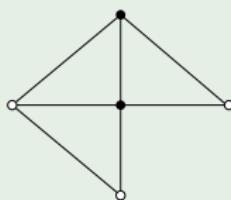


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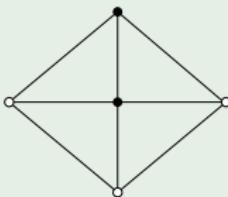


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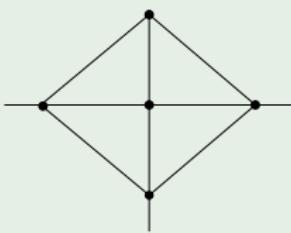


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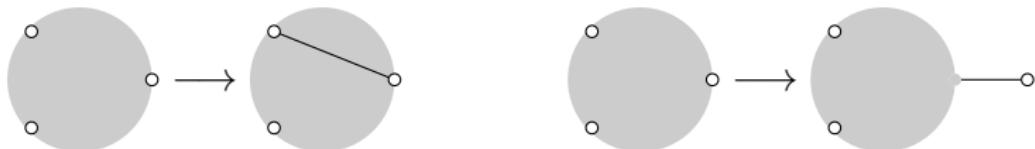


## Theorem (Panzer)

All  $\epsilon$ -coefficients of these graphs (off-shell) are MPL over the alphabet  $\{z, \bar{z}, 1 - z, 1 - \bar{z}, z - \bar{z}, 1 - z\bar{z}, 1 - z - \bar{z}, z\bar{z} - z - \bar{z}\}$ .

# Linearly reducible families (infinite)

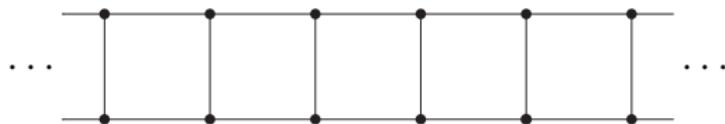
- 3-constructible graphs (3-point functions) [Brown, Schnetz, Panzer]



## Theorem (Panzer)

All  $\epsilon$ -coefficients of these graphs (off-shell) are MPL over the alphabet  $\{z, \bar{z}, 1 - z, 1 - \bar{z}, z - \bar{z}, 1 - z\bar{z}, 1 - z - \bar{z}, z\bar{z} - z - \bar{z}\}$ .

- minors of ladder-boxes (up to 2 legs off-shell)



## Theorem (Panzer)

All  $\epsilon$ -coefficients of these graphs are MPL. For the massless case, the alphabet is just  $\{x, 1 + x\}$  for  $x = s/t$ .